

Refilling Fixed-Capacity Containers

Gilles Gnacadja

Research and Development Information Systems, Amgen
Thousand Oaks, California, USA

<http://math.GillesGnacadja.info/>

California State University Channel Islands

Summer Mathematics

Research Experience for Undergraduates

18 July 2012

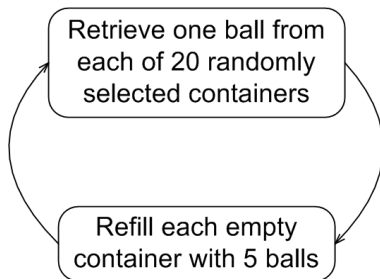
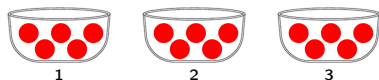
Outline

- 1 Steady-State Refilling in a Single-Stage Retrieval Process
 - The Single-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 2 Steady-State Refilling in a Multi-Stage Retrieval Process
 - The Multi-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 3 Asymptotic Equidistribution of Congruence Classes
 - The Theorem
 - The Proof

Outline

- 1 Steady-State Refilling in a Single-Stage Retrieval Process
 - The Single-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 2 Steady-State Refilling in a Multi-Stage Retrieval Process
 - The Multi-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 3 Asymptotic Equidistribution of Congruence Classes
 - The Theorem
 - The Proof

Retrieving and Refilling



How many containers are refilled at the n th pass?

$$n = 1$$

$$n = 2$$

$$n = 3$$

$$n = 4$$

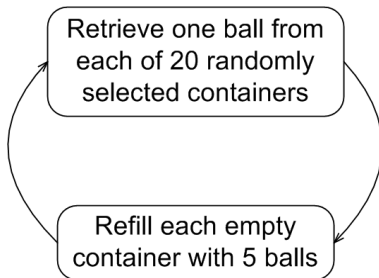
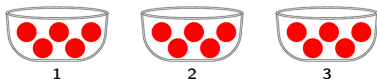
$$n = 5$$

$$n = 6$$

\vdots

$$n \rightarrow \infty$$

Retrieving and Refilling



How many containers are refilled at the n th pass?

$$n = 1 \quad \text{refilled} = 0$$

$$n = 2 \quad \text{refilled} = 0$$

$$n = 3 \quad \text{refilled} = 0$$

$$n = 4 \quad \text{refilled} = 0$$

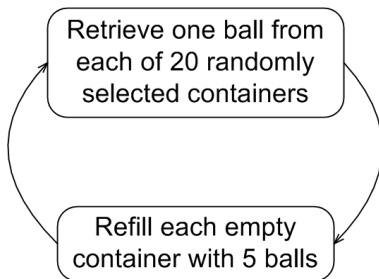
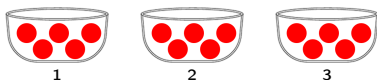
$$n = 5$$

$$n = 6$$

\vdots

$$n \rightarrow \infty$$

Retrieving and Refilling



How many containers are refilled at the n th pass?

$$n = 1 \quad \text{refilled} = 0$$

$$n = 2 \quad \text{refilled} = 0$$

$$n = 3 \quad \text{refilled} = 0$$

$$n = 4 \quad \text{refilled} = 0$$

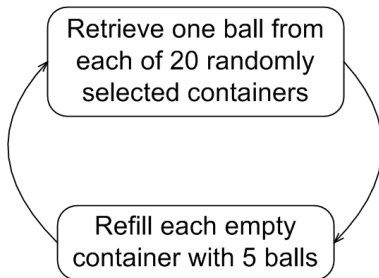
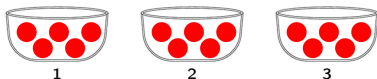
$$n = 5 \quad \text{refilled} = ?$$

$$n = 6 \quad \text{refilled} = ?$$

\vdots

$$n \rightarrow \infty$$

Retrieving and Refilling



How many containers are refilled at the n th pass?

$$n = 1 \quad \text{refilled} = 0$$

$$n = 2 \quad \text{refilled} = 0$$

$$n = 3 \quad \text{refilled} = 0$$

$$n = 4 \quad \text{refilled} = 0$$

$$n = 5 \quad \text{refilled} = ?$$

$$n = 6 \quad \text{refilled} = ?$$

\vdots

$$n \rightarrow \infty \quad \text{refilled} \rightarrow 4$$

Goal: Prove this

Outline

- 1 Steady-State Refilling in a Single-Stage Retrieval Process
 - The Single-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 2 Steady-State Refilling in a Multi-Stage Retrieval Process
 - The Multi-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 3 Asymptotic Equidistribution of Congruence Classes
 - The Theorem
 - The Proof

Refilling Events



Container i
is refilled at
the n th pass

=

The cumulative number of times container i has been selected since the first pass just became a multiple of $d = 5$ at the n th pass.

$\not\subseteq$

The cumulative number of times container i has been selected since the first pass is a multiple of $d = 5$ at the n th pass.

Refilling Events



Container i
is refilled at
the n th pass

=

The cumulative number of times container i has been selected since the first pass just became a multiple of d at the n th pass.

=

Container i
is selected
at the n th
pass

×

The cumulative number of times container i has been selected at passes $1, \dots, n-1$ is congruent to $d-1$ modulo d .

Probability of Refilling Events



Container i
 is refilled at
 the n th pass

=

Container i
 is selected
 at the n th
 pass

×

The cumulative number of
 times container i has been se-
 lected at passes $1, \dots, n - 1$ is
 congruent to $d - 1$ modulo d .



$$\text{probability} = \frac{M}{N}$$

($M := \#(\text{Selected})$)



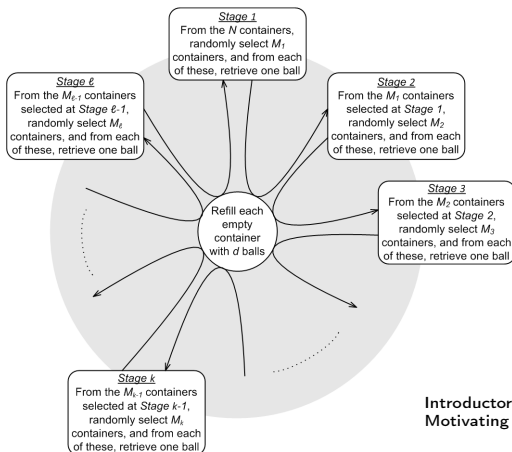
$$\text{probability} \xrightarrow{n \rightarrow \infty} \frac{1}{d}$$

(from upcoming theorem)

Outline

- 1 Steady-State Refilling in a Single-Stage Retrieval Process
 - The Single-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 2 Steady-State Refilling in a Multi-Stage Retrieval Process
 - The Multi-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 3 Asymptotic Equidistribution of Congruence Classes
 - The Theorem
 - The Proof

Multi-Stage Retrieving and Refilling



Introductory illustration: $l = 1$
Motivating application: $l = 2$

Outline

- 1 Steady-State Refilling in a Single-Stage Retrieval Process
 - The Single-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 2 Steady-State Refilling in a Multi-Stage Retrieval Process
 - The Multi-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 3 Asymptotic Equidistribution of Congruence Classes
 - The Theorem
 - The Proof

Probability of Refilling Events



Container i is refilled after stage k of the n th pass

=

Container i is selected at stages $1, \dots, k$ of the n th pass

×

The cumulative number of times container i has been selected at passes $1, \dots, n-1$ is congruent to $d - k$ modulo d .



probability = ...



probability $\xrightarrow{n \rightarrow \infty} 1/d$
(from upcoming theorem)

Outline

- 1 Steady-State Refilling in a Single-Stage Retrieval Process
 - The Single-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 2 Steady-State Refilling in a Multi-Stage Retrieval Process
 - The Multi-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 3 Asymptotic Equidistribution of Congruence Classes
 - The Theorem
 - The Proof

Probability Vector

$f(k)$ for $k = 0, \dots, \ell$: Probability that, in one pass, the number of times container i is selected is k

$f = (f(0), \dots, f(\ell))$: Probability vector, i.e.
 $f(0), \dots, f(\ell) \in \mathbb{R}_{\geq 0}$ and $f(0) + \dots + f(\ell) = 1$

Exercise: $f(k) = (M_k - M_{k+1})/N$ for $k = 0, \dots, \ell$,
with $M_0 := N$ and $M_{\ell+1} := 0$

Convolution of Probability Vectors

Probability that, in n passes,
the number of times container i is selected is k

n	Range of k	Probability
1	$0 \leq k \leq \ell$	$f(k)$
2	$0 \leq k \leq 2\ell$	$(f * f)(k) = f(0)f(k) + f(1)f(k-1) + f(2)f(k-2) + \dots$ $\dots + f(k-1)f(1) + f(k)f(0)$
\vdots		
n	$0 \leq k \leq n\ell$	$(\underbrace{f * f * \dots * f}_n)(k) = f^{*n}(k)$

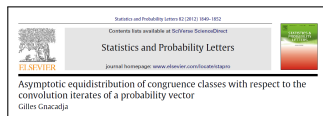
Convolution of vectors is like
multiplication of univariate polynomials.

Probability of Congruence Classes of Refilling Counts

Probability $\varphi(f, n, d, r)$ that, in n passes, the number of times container i is selected is congruent to r modulo d

$$\varphi(f, n, d, r) = \sum_{\substack{0 \leq k \leq nl \\ k \equiv r \pmod{d}}} f^{*n}(k) = \sum_{q=0}^{\text{floor}((nl-r)/d)} f^{*n}(r + dq)$$

Asymptotic Equidistribution of Congruence Classes



<http://dx.doi.org/10.1016/j.spl.2012.05.025>

Theorem

We have $\lim_{n \rightarrow \infty} \varphi(f, n, d, r) = 1/d$, where

- $f = (f(0), \dots, f(\ell))$ is a positive probability vector with $\ell \geq 1$;
- d and r are integers with $d \geq 1$ and $0 \leq r \leq d - 1$.

The message: *Congruence is asymptotically equidistributed regardless of (positive) starting distribution.*

Remark: Result may, but need not, fail if starting distribution f is nonnegative but not positive.

Outline

- 1 Steady-State Refilling in a Single-Stage Retrieval Process
 - The Single-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 2 Steady-State Refilling in a Multi-Stage Retrieval Process
 - The Multi-Stage Retrieval and Refilling Process
 - Mathematical Formulation of Refilling Events
- 3 Asymptotic Equidistribution of Congruence Classes
 - The Theorem
 - The Proof

Setting up for the Proof

Theorem to prove says that $\lim_{n \rightarrow \infty} \varphi(f, n, d)$ is the d vector filled with $1/d$, where

$$\varphi(f, n, d) := (\varphi(f, n, d, 0), \varphi(f, n, d, 1), \varphi(f, n, d, 2), \dots, \varphi(f, n, d, d-1)) .$$

$$\Phi(f, m, d) := \begin{pmatrix} \varphi(f, m, d, 0) & \varphi(f, m, d, 1) & \varphi(f, m, d, 2) & \cdots & \varphi(f, m, d, d-1) \\ \varphi(f, m, d, d-1) & \varphi(f, m, d, 0) & \varphi(f, m, d, 1) & \cdots & \varphi(f, m, d, d-2) \\ \varphi(f, m, d, d-2) & \varphi(f, m, d, d-1) & \varphi(f, m, d, 0) & \cdots & \varphi(f, m, d, d-3) \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \varphi(f, m, d, 1) & \varphi(f, m, d, 2) & \varphi(f, m, d, 3) & \cdots & \varphi(f, m, d, 0) \end{pmatrix}$$

$$= \text{CirculantMatrix}(\varphi(f, m, d))$$

- $\varphi(f, m, d)$ is a nonnegative probability vector; therefore...
- $\Phi(f, m, d)$ is a nonnegative doubly stochastic matrix.
- For m large enough (specifically $m \geq (d-1)/\ell$):
 - The probability vector $\varphi(f, m, d)$ is positive; therefore...
 - The doubly stochastic matrix $\Phi(f, m, d)$ is positive; therefore...
 - Perron-Frobenius Theory:
Powers of $\Phi(f, m, d)$ converge of the $d \times d$ matrix filled with $1/d$.

Elements of Proof

Proposition For $d, m, n \in \mathbb{Z}_{\geq 1}$ with $n > m$,
$$\varphi(f, n, d) = \varphi(f, n - m, d) \cdot \Phi(f, m, d).$$

Corollary For $d, m, s, t \in \mathbb{Z}_{\geq 1}$,
$$\varphi(f, sm + t, d) = \varphi(f, t, d) \cdot (\Phi(f, m, d))^s.$$

Consequence For $d, m, t \in \mathbb{Z}_{\geq 1}$ with $m \geq (d - 1)/\ell$,
$$\varphi(f, sm + t, d) \xrightarrow{s \rightarrow \infty} \underbrace{(1/d, \dots, 1/d)}_d.$$

Conclusion For $d \in \mathbb{Z}_{\geq 1}$, $\varphi(f, n, d) \xrightarrow{n \rightarrow \infty} \underbrace{(1/d, \dots, 1/d)}_d.$